

Lecture 2

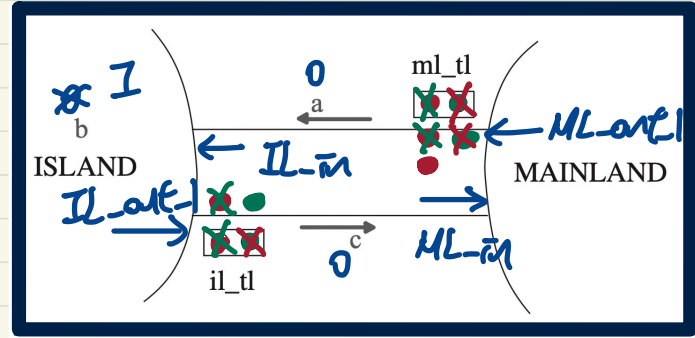
Part R

***Case Study on Reactive Systems -
Bridge Controller
2nd Refinement: Livelock/Divergence***

Current m2 May Livelock

ML_tl_green
when
 ✓ $ml_tl = red$
 ✓ $a + b < d$
 ✓ $c = 0$
then
 $ml_tl := green$
 $il_tl := red$
end

IL_tl_green
when
 $il_tl = red$
 $b > 0$
 $a = 0$
then
 $il_tl := green$
 $ml_tl := red$
end



$d=2$ Expected trace: no divergent fork turs trace

$\langle init, ML_tl_green, ML_out_1, IL_in, IL_tl_green, ML_tl_green, IL_tl_green, \dots \rangle$

a new event

old BOSS

Is ML_tl.g. enabled?

Is IL_tl.g. enabled?

→ also a valid trace of m2, but leading to livelock

\langle	$init$	ML_tl_green	ML_out_1	IL_in	IL_tl_green	ML_tl_green	IL_tl_green	$\dots \rangle$
	$d=2$	$d=2$	$d=2$	$d=2$	$d=2$	$d=2$	$d=2$	
	$a'=0$	$a'=0$	$a'=1$	$a'=0$	$a'=0$	$a'=0$	$a'=0$	
	$b'=0$	$b'=0$	$b'=0$	$b'=1$	$b'=1$	$b'=1$	$b'=1$	
	$c'=0$	$c'=0$	$c'=0$	$c'=0$	$c'=0$	$c'=0$	$c'=0$	
	$ml_tl = red$	$ml_tl' = green$	$ml_tl' = green$	$ml_tl' = green$	$ml_tl' = red$	$ml_tl' = green$	$ml_tl' = red$	
	$il_tl = red$	$il_tl' = red$	$il_tl' = red$	$il_tl' = red$	$il_tl' = green$	$il_tl' = red$	$il_tl' = green$	

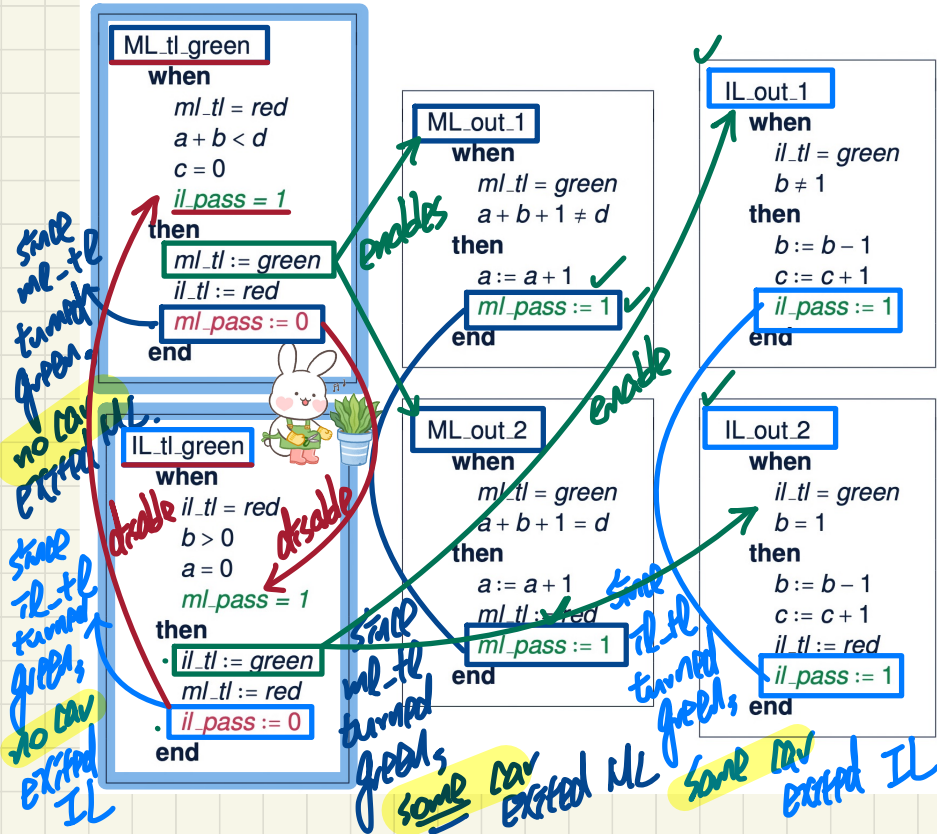


pattern of divergence

Fixing m2: Regulating Traffic Light Changes

To break the divergence pattern, after each view event occurring, some old events occur.

Divergence Trace: <init, ML_tl_green, ML_out_1, IL_in, IL_tl_green, ML_tl_green, IL_tl_green, ...>



d = 2	ml_pass	il_pass
< init,	1	1
ML_tl_green,	0	1
ML_out_1,	1	1
ML_out_2,	1	1
IL_in,	1	1
IL_in,	1	1
IL_tl_green,	1	0
IL_out_1,	1	1
IL_out_2,	1	1
ML_in,	1	1
ML_in	1	1
>		

Fixing m2: Measuring Traffic Light Changes

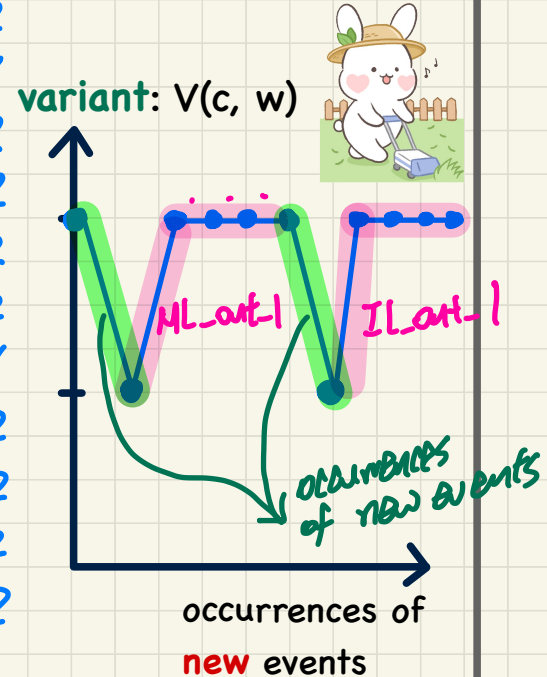
```

ML_tl_green
when
  ml_tl = red
  a + b < d
  c = 0
  il_pass = 1
then
  ml_tl := green
  il_tl := red
  ml_pass := 0
end
    
```

```

IL_tl_green
when
  il_tl = red
  b > 0
  a = 0
  ml_pass = 1
then
  il_tl := green
  ml_tl := red
  il_pass := 0
end
    
```

d = 2	ml_pass	il_pass	variants: <u>ml_pass + il_pass</u>
< init,	1	1	2
ML_tl_green,	0	1	1
ML_out_1,	1	1	2
ML_out_2,	1	1	2
IL_in,	1	1	2
IL_in,	1	1	2
IL_tl_green,	1	0	1
IL_out_1,	1	1	2
IL_out_2,	1	1	2
ML_in,	1	1	2
ML_in	1	1	2
>			



PO of Convergence/Non-Divergence/Livelock Freedom

A New Event Occurrence Decreases Variant

$$* \cancel{ml_pass^0} + \cancel{il_pass^{tl_pass}} < ml_pass + tl_pass$$

Variants: $ml_pass + il_pass$

ML_tl_green/VAR

$A(c)$
 $I(c, v)$
 $J(c, v, w)$
 $H(c, w)$
 \vdash
 $V(c, F(c, w)) < V(c, w)$

Post-state Evaluation
 Pre-state Evaluation

VAR
 applicable to new events

```

ML_tl_green
when
  ml_tl = red
  a + b < d
  c = 0
  il_pass = 1
then
  ml_tl := green
  il_tl := red
  ml_pass := 0
end
  
```

BAP:
 $ml_pass' = 0$
 $tl_pass' = tl_pass$



$d \in \mathbb{N}$	$d > 0$	
$COLOUR = \{green, red\}$	$green \neq red$	
$n \in \mathbb{N}$	$n \leq d$] m0
$a \in \mathbb{N}$	$b \in \mathbb{N}$	
$a + b + c = n$	$a = 0 \vee c = 0$] m1
$ml_tl \in COLOUR$	$il_tl \in COLOUR$	
$ml_tl = green \Rightarrow a + b < d \wedge c = 0$	$il_tl = green \Rightarrow b > 0 \wedge a = 0$] m2
$ml_tl = red \vee il_tl = red$		
$ml_pass \in \{0, 1\}$	$il_pass \in \{0, 1\}$] m3
$ml_tl = red \Rightarrow ml_pass = 1$	$il_tl = red \Rightarrow il_pass = 1$	
$ml_tl = red$	$a + b < d$] c = 0
$il_pass = 1$		

\vdash
 $0 + il_pass < ml_pass + il_pass$

Concrete guards of ML_tl_green

Lecture 2

Part S

***Case Study on Reactive Systems -
Bridge Controller
2nd Refinement:
Relative Deadlock Freedom***

PO of Relative Deadlock Freedom

```

axm0.1  d ∈ ℕ
axm0.2  d > 0
axm2.1  COLOUR = {green, red}
axm2.2  green ≠ red
inv0.1  n ∈ ℕ
inv0.2  n ≤ d
inv1.1  a ∈ ℕ
inv1.2  b ∈ ℕ
inv1.3  c ∈ ℕ
inv1.4  a + b + c = n
inv1.5  a = 0 ∨ c = 0
inv2.1  ml_tl ∈ COLOUR
inv2.2  il_tl ∈ COLOUR
inv2.3  ml_tl = green ⇒ a + b < d ∧ c = 0
inv2.4  il_tl = green ⇒ b > 0 ∧ a = 0
inv2.5  ml_tl = red ∨ il_tl = red
inv2.6  ml_pass ∈ {0, 1}
inv2.7  il_pass ∈ {0, 1}
inv2.8  ml_tl = red ⇒ ml_pass = 1
inv2.9  il_tl = red ⇒ il_pass = 1
    
```

Disjunction of abstract guards



Disjunction of concrete guards

guards of ML.out in m_1
 $a + b < d \wedge c = 0$
 $c > 0$
 $a > 0$
 $b > 0 \wedge a = 0$

guards of ML.in in m_1
 $a > 0$

guards of IL.in in m_1
 $a > 0$

guards of IL.out in m_1
 $b > 0 \wedge a = 0$

guards of ML_tl.green in m_2
 $ml_tl = red \wedge a + b < d \wedge c = 0 \wedge il_pass = 1$

guards of IL_tl.green in m_2
 $il_tl = red \wedge b > 0 \wedge a = 0 \wedge ml_pass = 1$

guards of ML.out.1 in m_2
 $ml_tl = green \wedge a + b + 1 \neq d$

guards of ML.out.2 in m_2
 $ml_tl = green \wedge a + b + 1 = d$

guards of IL.out.1 in m_2
 $il_tl = green \wedge b \neq 1$

guards of IL.out.2 in m_2
 $il_tl = green \wedge b = 1$

guards of ML.in in m_2
 $a > 0$

guards of IL.in in m_2
 $c > 0$

Abstract m_1

variables: a, b, c

invariants:

```

inv1.1 : a ∈ ℕ
inv1.2 : b ∈ ℕ
inv1.3 : c ∈ ℕ
inv1.4 : a + b + c = n
inv1.5 : a = 0 ∨ c = 0
    
```

ML.out

when

$a + b < d$
 $c = 0$

then

$a := a + 1$

end

ML.in

when

$c > 0$

then

$c := c - 1$

end

IL.in

when

$a > 0$

then

$a := a - 1$

$b := b + 1$

end

IL.out

when

$b > 0$

$a = 0$

then

$b := b - 1$

$c := c + 1$

end

Concrete m_2

ML_tl.green

when

$ml_tl = red$

$a + b < d$

$c = 0$

$il_pass = 1$

then

$ml_tl := green$

$il_tl := red$

$ml_pass := 0$

end

IL_tl.green

when

$il_tl = red$

$b > 0$

$a = 0$

$ml_pass = 1$

then

$il_tl := green$

$ml_tl := red$

$il_pass := 0$

end

ML.out.1

when

$ml_tl = green$

$a + b + 1 \neq d$

then

$a := a + 1$

$ml_pass := 1$

end

IL.out.1

when

$il_tl = green$

$b \neq 1$

then

$b := b - 1$

$c := c + 1$

$il_pass := 1$

end

ML.out.2

when

$ml_tl = green$

$a + b + 1 = d$

then

$a := a + 1$

$ml_tl := red$

$ml_pass := 1$

end

IL.out.2

when

$il_tl = green$

$b = 1$

then

$b := b - 1$

$c := c + 1$

$il_tl := red$

$il_pass := 1$

end

IL.in

when

$a > 0$

then

$a := a - 1$

$b := b + 1$

end

ML.in

when

$c > 0$

then

$c := c - 1$

end

Discharging **POs** of m2: **Relative Deadlock Freedom**

```

d ∈ ℕ
d > 0
COLOUR = {green, red}
green ≠ red
n ∈ ℕ
n ≤ d
a ∈ ℕ
b ∈ ℕ
c ∈ ℕ
a + b + c = n
a = 0 ∨ c = 0
ml_tl ∈ COLOUR
il_tl ∈ COLOUR
ml_tl = green ⇒ a + b < d ∧ c = 0
il_tl = green ⇒ b > 0 ∧ a = 0
ml_tl = red ∨ il_tl = red
ml_pass ∈ {0, 1}
il_pass ∈ {0, 1}
ml_tl = red ⇒ ml_pass = 1
il_tl = red ⇒ il_pass = 1
  a + b < d ∧ c = 0
  ∨ c > 0
  ∨ a > 0
  ∨ b > 0 ∧ a = 0
┌
  ml_tl = red ∧ a + b < d ∧ c = 0 ∧ il_pass = 1
  ∨ il_tl = red ∧ b > 0 ∧ a = 0 ∧ ml_pass = 1
  ∨ ml_tl = green
  ∨ il_tl = green
  ∨ a > 0
  ∨ c > 0
    
```



Study

Ex. 1

⋮

```

d ∈ ℕ
d > 0
b ∈ ℕ
ml_tl = red
il_tl = red
ml_tl = red ⇒ ml_pass = 1
il_tl = red ⇒ il_pass = 1
┌
  b < d ∧ ml_pass = 1 ∧ il_pass = 1
  ∨ b > 0 ∧ ml_pass = 1 ∧ il_pass = 1
    
```

Ex. 2

⋮

```

d ∈ ℕ
d > 0
b ∈ ℕ
ml_tl = red
il_tl = red
ml_pass = 1
il_pass = 1
┌
  b < d ∧ ml_pass = 1 ∧ il_pass = 1
  ∨ b > 0 ∧ ml_pass = 1 ∧ il_pass = 1
    
```

Ex. 3

⋮

```

d > 0
b ∈ ℕ
┌
  b < d ∨ b > 0
    
```

ARI

```

d > 0
b > 0 ∨ b = 0
┌
  b < d ∨ b > 0
    
```

OR.L

```

d > 0
b > 0
┌
  b < d ∨ b > 0
    
```

OR.R2

```

d > 0
b > 0
┌
  b > 0
    
```

HYP

```

d > 0
b = 0
┌
  b < d ∨ b > 0
    
```

EQ.LR, MON

```

d > 0
┌
  0 < d ∨ 0 > 0
    
```

OR.R1

```

d > 0
┌
  0 < d
    
```

HYP

1st Refinement and 2nd Refinement: Provably Correct

Abstract m1

variables: a, b, c

constants: d

axioms:
 $axm0.1: d \in \mathbb{N}$
 $axm0.2: d > 0$

invariants:
 $inv1.1: a \in \mathbb{N}$
 $inv1.2: b \in \mathbb{N}$
 $inv1.3: c \in \mathbb{N}$
 $inv1.4: a + b + c = n$
 $inv1.5: a = 0 \vee c = 0$

variants:
 $2 \cdot a + b$

init
begin
 $a := 0$
 $b := 0$
 $c := 0$
end

ML.out
when
 $a + b < d$
 $c = 0$
then
 $a := a + 1$
end

ML.in
when
 $c > 0$
then
 $c := c - 1$
end

IL.in
when
 $a > 0$
then
 $a := a - 1$
 $b := b + 1$
end

IL.out
when
 $b > 0$
 $a = 0$
then
 $b := b - 1$
 $c := c + 1$
end

Correctness Criteria:

- + Guard Strengthening
- + Invariant Establishment
- + Invariant Preservation
- + Convergence
- + Relative Deadlock Freedom



Art

superposition

variables:
 a
 b
 c
 ml_tl
 il_tl
 ml_pass
 il_pass

constants: d

sets: $COLOR$

axioms:
 $axm0.1: d \in \mathbb{N}$
 $axm0.2: d > 0$
 $axm2.1: COLOR = \{green, red\}$
 $axm2.2: green \neq red$

invariants:
 $inv2.1: ml_tl \in COLOR$
 $inv2.2: il_tl \in COLOR$
 $inv2.3: ml_tl = green \Rightarrow a + b < d \wedge c = 0$
 $inv2.4: il_tl = green \Rightarrow b > 0 \wedge a = 0$
 $inv2.5: ml_tl = red \vee il_tl = red$
 $inv2.6: ml_pass \in \{0, 1\}$
 $inv2.7: il_pass \in \{0, 1\}$
 $inv2.8: ml_tl = red \Rightarrow ml_pass = 1$
 $inv2.9: il_tl = red \Rightarrow il_pass = 1$

variants:
 $ml_pass + il_pass$

ML.tl.green
when
 $ml_tl = red$
 $a + b < d$
 $c = 0$
 $il_pass = 1$
then
 $ml_tl := green$
 $il_tl := red$
 $ml_pass := 0$
end

IL.tl.green
when
 $il_tl = red$
 $b > 0$
 $a = 0$
 $ml_pass = 1$
then
 $il_tl := green$
 $ml_tl := red$
 $il_pass := 0$
end

ML.out.1
when
 $il_tl = green$
 $a + b + 1 \neq d$
then
 $a := a + 1$
 $ml_pass := 1$
end

IL.out.1
when
 $il_tl = green$
 $b \neq 1$
then
 $b := b - 1$
 $c := c + 1$
 $il_pass := 1$
end

ML.in
when
 $c > 0$
then
 $c := c - 1$
end

ML.out.2
when
 $ml_tl = green$
 $a + b + 1 = d$
then
 $a := a + 1$
 $ml_tl := red$
 $ml_pass := 1$
end

IL.out.2
when
 $il_tl = green$
 $b = 1$
then
 $b := b - 1$
 $c := c + 1$
 $il_tl := red$
 $il_pass := 1$
end

IL.in
when
 $a > 0$
then
 $a := a - 1$
 $b := b + 1$
end

disjoint freedom

Concrete m2